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# Computing Testing Equivalence with Binary Decision Diagrams

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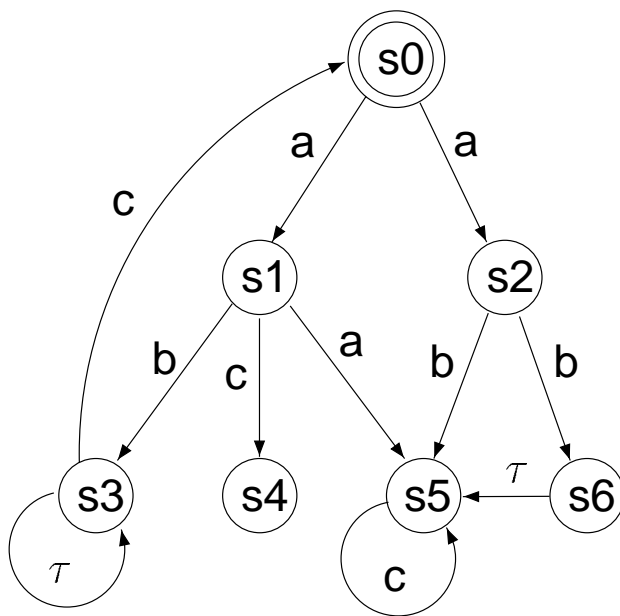




## Introduction

- **Process algebras** are a convenient tool for describing and reasoning about the behaviour of concurrent systems.
- Systems are described as transition graphs called **processes**.
- **Equivalence relations** are used to study the relationship between different systems or between different levels of abstractions of the system.

## Process



$$P = (\mathcal{S}, Act, \delta, p_0)$$

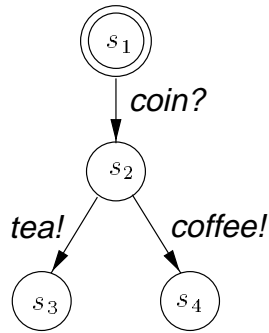
- $\mathcal{S}$ : set of states
- $Act$ : set of actions
- $\delta$ : transition relation
- $p_0$ : initial state

- $Act$  contains **visible actions** and **silent action**  $\tau$ .
- $p$  is a **deadlock state** iff no actions can be performed from it.
- $p$  is a **divergent state** iff there is an infinite sequence of silent actions starting in it.

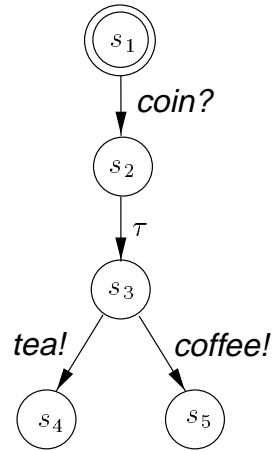




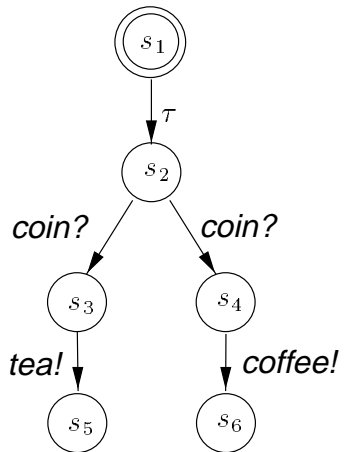
# Some different descriptions of drink machine



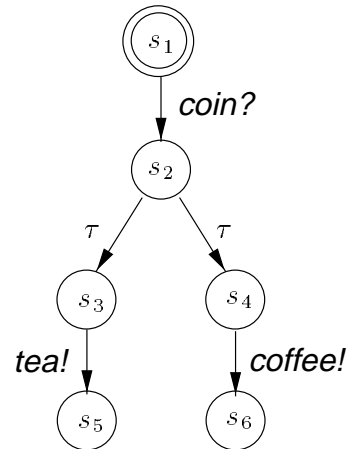
Drink machine A1



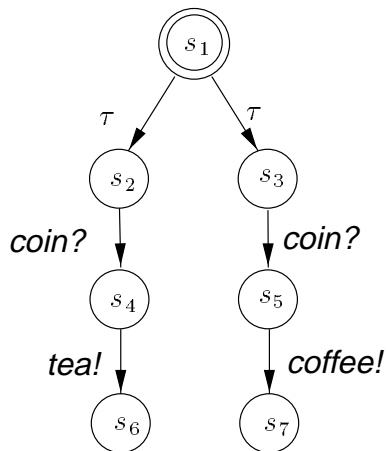
Drink machine A2



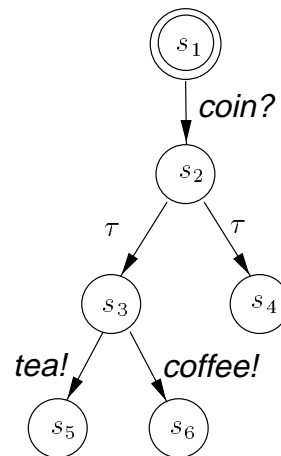
Drink machine A3



Drink machine A4



Drink machine A5



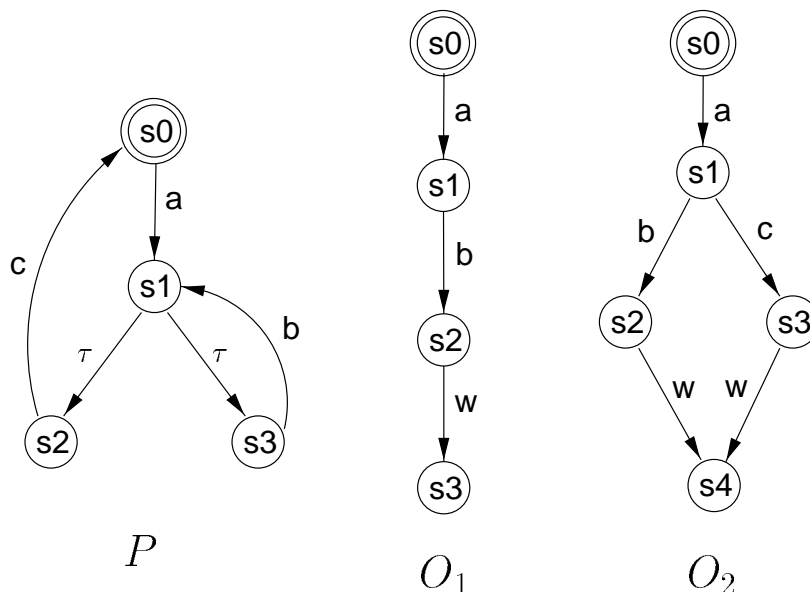
Drink machine A6





## Definition of Testing Equivalence

- An **observer** is a process which runs concurrently with the observed process and interacts with it. The observer uses special action  $w$  to report to the external world the **success of the observation**.
- The process  $P$  **must satisfy** the observer  $O$  if during the observation the action  $w$  is inevitable. The process  $P$  **may satisfy** the observer  $O$  if during the observation the action  $w$  is possible.
- Two processes are **must equivalent** iff they must satisfy the same sets of observers. They are **may equivalent** iff they may satisfy the same sets of observers. The processes are **testing equivalent** iff they are must equivalent and may equivalent.





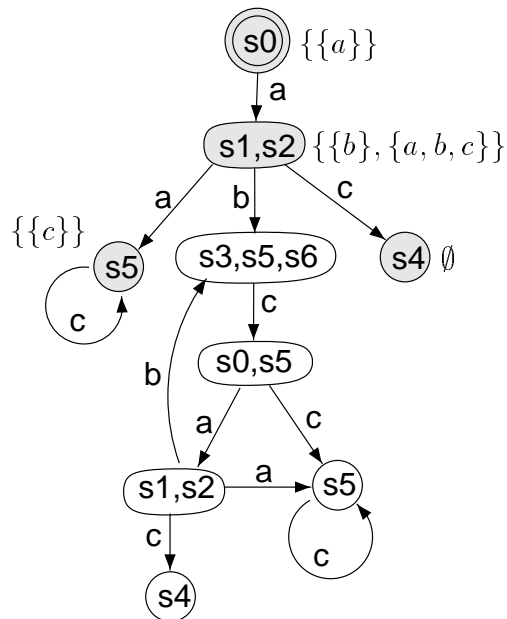
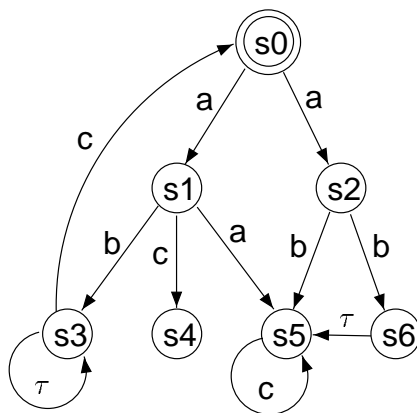
# Testing Equivalence via Bisimulation

The acceptance graph  $P'$  is an annotated deterministic process obtained from process  $P$  as follows:

- Every state in  $P'$  is a **set of states** of process  $P$ .
- A **visible action**  $\alpha$  is possible in the state in  $P'$  iff the same action is possible in one of matched states in  $P$ .
- The states in  $P'$  can be **open** or **closed**. Closed states are annotated with a minimized **acceptance sets**.

The **acceptance set** of a closed state in  $P'$  is a set of actions that are immediately possible from matched states in  $P$ .

Two processes are **testing equivalent** if their acceptance graphs are bisimulation equivalent.





## Conclusion

- **Testing equivalence** is conceptually simple and close to the concept of testing.
- We made **efficient implementation** of a testing equivalence using Binary Decision Diagrams (BDDs).
- **The algorithm** for testing equivalence is composed of transforming processes into acceptance graphs and then determining bisimulation equivalence between them.
- **Possible extensions:**
  - computing testing preorders,
  - system abstraction due to testing equivalence,
  - generating diagnostics to explain why two systems are not testing equivalent.
- Presented work is a part of **our verification system** also capable of:
  - changing the behaviour of processes,
  - performing parallel composition of processes,
  - determining bisimulation equivalence between processes,
  - reasoning about properties of processes using model checking with ACTL.

